# Lecture 10, 17 September 2024

#### Arrays

- Contiguous block of memory
- Typically size is declared in advance, all values are uniform
- a[0] points to first memory location in the allocated block
- Locate a[i] in memory using index arithmetic
  - Skip i blocks of memory, each block's size determined by value stored in array
- Random access -- accessing the value at a[i] does not depend on i
- Useful for procedures like sorting, where we need to swap out of order values a[i] and a[j]
  - a[i], a[j] = a[j], a[i]
  - Cost of such a swap is constant, independent of where the elements to be swapped are in the array
- Inserting or deleting a value is expensive
- · Need to shift elements right or left, respectively, depending on the location of the modification

## Lists

- Each location is a *cell*, consisiting of a value and a link to the next cell
- Think of a list as a train, made up of a linked sequence of cells
- The name of the list 1 gives us access to 1[0], the first cell
- To reach cell [[i], we must traverse the links from 1[0] to 1[1] to 1[2] ... to 1[i-1]] to 1[i]
  - Takes time proportional to i
- Cost of swapping l[i] and l[j] varies, depending on values i and j
- On the other hand, if we are already at l[i] modifying the list is easy
  - Insert create a new cell and reroute the links
  - Delete bypass the deleted cell by rerouting the links
- · Each insert/delete requires a fixed amount of local "plumbing", independent of where in the list it is performed

## Dictionaries

- Values are stored in a fixed block of size m
- Keys are mapped to  $\{0, 1, ..., m 1\}$
- Hash function h: K 
  ightarrow S maps a  $\mathit{large}$  set of keys K to a  $\mathit{small}$  range S
- Simple hash function: interpret  $k \in K$  as a bit sequence representing a number  $n_k$  in binary, and compute  $n_k \mod m$ , where |S| = m
- Mismatch in sizes means that there will be *collisions* --  $k_1 \neq k_2$ , but  $h(k_1) = h(k_2)$
- A good hash function maps keys "randomly" to minimize collisions
- Hash can be used as a *signature* of authenticity
- Modifying k slightly will drastically alter h(k)
  - No easy way to reverse engineer a k' to map to a given h(k)
  - Use to check that large files have not been tampered with in transit, either due to network errors or malicious intervention
- Dictionary uses a hash function to map key values to storage locations
- Lookup requires computing h(k) which takes roughly the same time for any k
  - Compare with computing the offset a[i] for any index i in an array
- · Collisions are inevitable, different mechanisms to manage this, which we will not discuss now
- Effectively, a dictionary combines flexibility with random access

# Lists in Python

- Flexible size, allow inserting/deleting elements in between
- · However, implementation is an array, rather than a list
- · Initially allocate a block of storage to the list
- When storage runs out, double the allocation
- l.append(x) is efficient, moves the right end of the list one position forward within the array
- l.insert(0,x) inserts a value at the start, expensive because it requires shifting all the elements by 1
- We will run experiments to validate these claims

#### Measuring execution time

- Call time.perf\_counter()
- Actual return value is meaningless, but difference between two calls measures time in seconds

•  $10^7$  appends to an empty Python list

• Doubling the work approximately doubles the time, linear

+  $10^5$  inserts at the beginning of a Python list

- Doubling and tripling the work multiplies the time by 4 and 9, respectively, so quadratic

```
11.383465348975733
```

• Creating  $10^7$  entries in an empty dictionary

```
In [7]: start = time.perf_counter()
d = {}
for i in range(10000000,0,-1):
    d[i] = i
elapsed = time.perf_counter() - start
print(elapsed)
```

```
1.0553472369792871
```

- Doubling the operations, doubles the time, so linear
- Dictionaries are effectively random access

```
In [9]: start = time.perf_counter()
d = {}
for i in range(20000000,0,-1):
    d[i] = i
    elapsed = time.perf_counter() - start
    print(elapsed)
2.602921769954264
```

- Insert keys in random order
- Use the library function random.shuffle(l) to permute the elements of l

```
In [10]: import random
```

lhundred = list(range(100))
random.shuffle(lhundred)
print(lhundred)

[11, 37, 34, 12, 46, 41, 96, 6, 16, 13, 97, 76, 26, 47, 27, 28, 99, 62, 90, 0, 51, 81, 79, 35, 5, 48, 84, 53, 6 5, 85, 25, 82, 52, 57, 78, 23, 98, 54, 20, 63, 91, 19, 38, 75, 80, 7, 3, 64, 74, 2, 31, 72, 93, 39, 56, 71, 14, 30, 77, 40, 55, 43, 68, 69, 61, 29, 33, 9, 44, 36, 15, 32, 18, 94, 21, 24, 60, 49, 70, 22, 45, 92, 89, 17, 58, 1 0, 73, 66, 50, 59, 87, 4, 8, 1, 95, 88, 83, 67, 42, 86]

- Insert  $10^6$  keys in random order
- Note that we start the counter after we shuffle the list of keys, so we count only the time required to populate the dictionary

```
In [11]: import random
    keylist = list(range(1000000,0,-1))
```

```
rndkeylist = keylist[:]
random.shuffle(rndkeylist)
d = {}
start = time.perf_counter()
for i in keylist:
    d[i] = i
elapsed = time.perf_counter() - start
print("Sequential keys:", elapsed)
```

```
d = {}
start = time.perf_counter()
for i in rndkeylist:
    d[i] = i
elapsed = time.perf_counter() - start
print("Shuffled keys:", elapsed)
```

Sequential keys: 0.09673382097389549 Shuffled keys: 0.39740611804882064

• Double the number of keys to  $2 imes 10^6$ 

```
In [12]: import random
    keylist = list(range(2000000,0,-1))
    rndkeylist = keylist[:]
    random.shuffle(rndkeylist)
```

```
d = {}
start = time.perf_counter()
for i in keylist:
    d[i] = i
elapsed = time.perf_counter() - start
print("Sequential keys:", elapsed)
```

```
d = {}
start = time.perf_counter()
for i in rndkeylist:
    d[i] = i
elapsed = time.perf_counter() - start
print("Shuffled keys:", elapsed)
```

Sequential keys: 0.21819286403479055 Shuffled keys: 0.6841557070147246

• Triple the number of keys to  $3 imes 10^6$ 

```
In [14]: import random
         keylist = list(range(3000000,0,-1))
         rndkeylist = keylist[:]
         random.shuffle(rndkeylist)
         d = \{\}
         start = time.perf_counter()
         for i in keylist:
            d[i] = i
         elapsed = time.perf_counter() - start
         print("Sequential keys:", elapsed)
         d = {}
         start = time.perf_counter()
         for i in rndkeylist:
             d[i] = i
         elapsed = time.perf counter() - start
         print("Shuffled keys:", elapsed)
        Sequential keys: 0.35756950796348974
        Shuffled keys: 1.1829602149664424
```

- Using shuffled keys is about 3 times slower than inserting keys in sequence
- However, even after shuffling, the time taken grows approximately linearly

Implementing a "real" list using dictionaries

```
In [15]: def createlist(): # Equivalent of l = [] is l = createlist()
           return({})
         def listappend(l,x):
           if l == {}:
             l["value"] = x
             l["next"] = {}
             return
           node = l
           while node["next"] != {}:
            node = node["next"]
           node["next"]["value"] = x
           node["next"]["next"] = {}
           return
         def listinsert(l,x):
           if l == {}:
             l["value"] = x
             l["next"] = {}
             return
           newnode = {}
           newnode["value"] = l["value"]
           newnode["next"] = l["next"]
           l["value"] = x
           l["next"] = newnode
           return
         def printlist(l):
           print("{",end="")
           if l == {}:
            print("}")
             return
           node = l
           print(node["value"],end="")
           while node["next"] != {}:
             node = node["next"]
           print(",", node["value"], end="")
print("}")
           return
```

• Display a small list as nested dictionaries

```
In [16]: start = time.perf_counter()
    l = createlist()
    for i in range(10):
        listappend(l,i)
    elapsed = time.perf_counter() - start
    print(elapsed)
    print(l)
```

```
0.013133806001860648
{'value': 0, 'next': {'value': 1, 'next': {'value': 2, 'next': {'value': 3, 'next': {'value': 4, 'next': {'value': 5, 'next': {'value': 6, 'next': {'value': 7, 'next': {'value': 8, 'next': {'value': 9, 'next': {}}}}}}
```

- Insert  $10^7$  elements at the beginning in this implementation of a list

```
In [21]: start = time.perf_counter()
    l = createlist()
    for i in range(1000000):
        listinsert(l,i)
    elapsed = time.perf_counter() - start
    print(elapsed)
    1.2849651229917072
```

• Doubling the work doubles the time, so linear

```
for i in range(2000000):
    listinsert(l,i)
elapsed = time.perf_counter() - start
print(elapsed)
```

3.5748096029856242

- Append  $10^4$  elements in this implementation of a list

```
In [23]: start = time.perf_counter()
    l = createlist()
    for i in range(10000):
        listappend(l,i)
    elapsed = time.perf_counter() - start
    print(elapsed)
```

2.831144590047188

• Halving the work takes 1/4 of the time, so quadratic

```
In [24]: start = time.perf_counter()
    l = createlist()
    for i in range(5000):
        listappend(l,i)
    elapsed = time.perf_counter() - start
print(elapsed)
```

0.6491393339820206

## Defining our own data structures

- We have implemented a "linked" list using dictionaries
- The fundamental functions like listappend , listinsert , listdelete modify the underlying list
- Instead of mylist = {}, we wrote mylist = createlist()
- To check empty list, use a function isempty() rather than mylist == {}
- Can we clearly separate the interface from the implementation
- Define the data structure in a more "modular" way